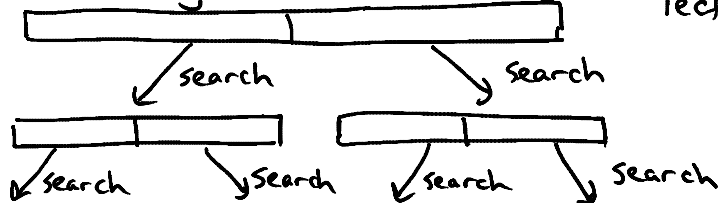


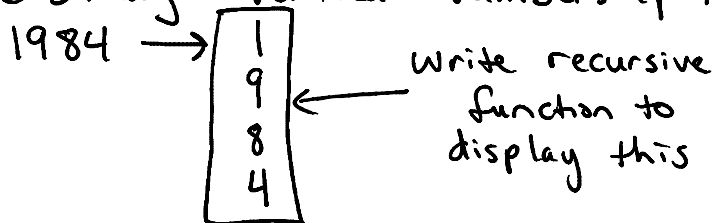
14.1 Recursive Functions for Tasks

recursion - a function that calls itself

Ex: binary search (p 785, cover next lecture)



Case Study - Vertical Numbers (p 765)



Think about how to break down problem
if number is one digit (< 10)
just print to screen
else

1) print 1 to $n-1$ digits to screen

2) print n th digit to screen

(1) leads to recursive function

call function w/ n th digit removed
number / 10 removes last digit

when using integers

(2) n th digit is $\text{number} \% 10$

This leads to pseudocode of alg
if number < 10 // stopping condition
print number

else

write-vertical (number / 10)

print number % 10

This has good recursive structure
one branch has no recursive call
stopping condition to prevent

endless function calls
one branch has recursive call
that breaks task into subset

Code

```
void write-vertical(int n);  
void write-vertical(int n)  
{  
    if (n < 10)  
    {  
        printf("%d\n", n);  
    }  
    else  
    {  
        write-vertical(n/10);  
        printf("%d\n", n%10);  
    }  
}
```

Tracing a Function Call

Book uses write-vertical(123) on p767

Use 1984 in class

write-vertical(1984)

if(1984 < 10) F

else

write-vertical(1984/10)

if(198 < 10) F

else

write-vertical(198/10)

if(19 < 10) F

else

write-vertical(19/10)

if(1 < 10) T

output 1,

output 19%10,

output 198%10,

output 1984%10,

1
9
8
4

on screen

output 198 \rightarrow 10,

9
4

output 1984 \rightarrow 10,

How It Works

When any function call reached, system calls function & waits for results before continuing to next line

Saves all info needed to continue when called function finishes

So when write-vertical (1984) calls write-vertical (198), it must wait for that to finish completely before continuing to the printf line

Stopping Recursion

The recursive calls have to stop at some point to allow all to finish

Have "base case" or "stopping case"

does not have a recursive call

Every possible input must eventually decompose to stopping case or some calls will never complete

Infinite Recursion - Pitfall

stopping case not reached or no stopping case at all

will typically run until resources exhausted or computer terminates program / crashes

if observed in code, check that

computations correctly lead to stopping case

Stacks for Recursion

a special memory structure used for function calls, organized LIFO (last in, first out)

only data at top of stack is readable because it is LIFO

push - put data on stack

pop - take top data off stack

how it is used

before a function call is executed,
certain information about currently
executing code is saved to the stack
- anything needed to continue execution
after function call exits

as function is executing, if it has
a function call it also pushes
data onto stack

when function call completes, results
noted & saved info about caller
is removed from stack

caller continues execution

activation frame - portion of memory that
contains info about caller

Stack Overflow

too many activation frames pushed
exceeds allocated max mem of stack

Ex: infinite recursion causes stack overflow

Recursion vs Iteration

Iterative version of write_vertical (p 777)

```
void write_vertical (int n)
```

```
{
```

```
    int tens = 1;
```

```
    int left_end = n;
```

```
    while (left_end > 10)
```

```
    {
```

```
        left_end = left_end / 10;
```

```
        tens = tens * 10; // find highest multiplier
```

```
    }
```

```
    // print from highest multiplier (1st digit)
```

```
    // to last digit
```

```
    for (int mod = tens; mod > 0; mod = mod / 10)
```

```
    {
```

```
        printf("%0d\n", n / mod);
```

```

    ' n = n % mod; '
  }
}

```

Iterative version not always as easy to come up with as recursive version

Recursive is not as efficient due to having to save context on stack

- often takes more memory because of stack
- doing a function call always has costs

Weigh benefits of simpler alg vs costs of stack

14.2 Recursive Functions for Values

Similar to void functions shown in 14.1 but returns a value

Many cases for using recursion to calculate

$$\text{factorial}(n) = n * \text{factorial}(n-1)$$

$$\text{power}(n, x) = n * \text{power}(n, x-1)$$

Base case should return base value

factorial(n) returns 1 when $n == 1$

power(n, x) returns 1 when $x == 0$

Code

```

int factorial(int n)
{
    if (n <= 1) return 1;
    else return n * factorial(n-1);
}

// Power function from p 780
int power(int base, int pow)
{
    if (pow < 0)
    {
        printf("Illegal power %d\n", pow);
        exit(1);
    }
    if (pow > 0) return power(base, pow-1);
}

```

```

    {
    if (pow > 0) return power (base, pow-1);
    else return 1;
    }

```

Tracing a factorial call

factorial (4)

if (4 <= 1) F

else return 4 * factorial (3)

if (3 <= 1) F

else return 3 * factorial (2)

if (2 <= 1) F

else return 2 * factorial (1)

if (1 <= 1) T

return 1

$2 * 1 = 2$

$3 * 2 = 6$

$4 * 6 = 24$

14.3 Thinking Recursively

Design Techniques

Check that chain of calls reaches stopping

case & stopping case is correct

Don't need to trace all calls like above

Instead, check following:

1) No infinite recursion

2) Stopping case(s) are correct

3) Final returned value is correct

(or action is correct for void funcs)

Case Study - Binary Search

Efficient way to search sorted array

Array must be sorted for this to work

This code will deviate from book's example

Book: two call-by-ref params

one bool - true if value found

one int - index of element if found

one int - index of element if found
Lecture: function will return int
if $int \geq 0$, index of found element
if $int == -1$, element not found

Pseudocode

if no elements, return NOT_FOUND
compute midpoint between start & end
if midpoint is element, return index
if element < midpoint
 search start to midpoint - 1
else element > midpoint
 search midpoint + 1 to end

Code

```
int binary-search (int a[],  
                  int start, int end, int elem)  
{  
    int mid;  
    if (start > end) return -1;  
    else  
    {  
        mid = (start + end) / 2;  
        if (elem == a[mid])  
            return mid;  
        else if (elem < a[mid])  
            return binary-search(a,  
                                start, mid - 1, elem);  
        else  
            return binary-search(a,  
                                mid + 1, end, elem);  
    }  
}
```

Checking the Recursion

- 1) No infinite recursion
 returns -1 when no more elements
 always reducing num of elements
- 2) Stopping case correct

2 stopping cases: not found & found
not found - no elements so nothing
could be found, correct
found - returns index where found

3) Final value is correct
-1 when not found
index when found

Programming Example - Recursive Member Func
member functions can use recursion too
coded the same way as stand-alone funcs
Example is to add x years of interest to
the balance of a bank account

Pseudocode

```
if years == 1
    add 1 yr of interest
if years > 1
    add 1 yr of interest
    call function w/ year - 1
```

Code

```
void BankAccount::update()
{
    balance += fraction(interest_rate) * balance;
}
void BankAccount::update(int years)
{
    if (years == 1) update();
    else if (years > 1)
    {
        update();
        update(years - 1);
    }
}
```